

# A NOTE ON DOMINATING SETS AND AVERAGE DISTANCE

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ABSTRACT. We show that the total domination number of a simple connected graph is greater than the average distance of the graph minus one-half, and that this inequality is best possible. In addition, we show that the domination number of the graph is greater than two-thirds of the average distance minus one-third, and that this inequality is best possible. Although the latter inequality is a corollary to a result of P. Dankelmann, we present a short and direct proof.

## 1. INTRODUCTION AND KEY DEFINITIONS

Let  $G = (V, E)$  be a simple connected graph of finite order  $|V| = n$ . Although we may identify a graph  $G$  with its set of vertices, in cases where we need to be explicit we write  $V(G)$  to denote the vertex set of  $G$ . A set  $D$  of vertices of a graph  $G$  is called a *dominating set* provided each vertex of  $V - D$  is adjacent to a member of  $D$ . The *domination number* of  $G$ , denoted  $\gamma(G)$ , is the cardinality of a smallest dominating set in  $G$ . Likewise, a set  $D$  of vertices is called a *total dominating set* provided each vertex of  $V$  is adjacent to a member of  $D$ . The *total domination number* of  $G$ , denoted  $\gamma_t(G)$ , is the cardinality of a smallest total dominating set in  $G$ . The *distance* between two vertices  $u$  and  $v$  in  $G$  is the length of a shortest path in  $G$  connecting  $u$  and  $v$ . The *Wiener index* or *total distance* of  $G$ , denoted by  $W = W(G)$ , is the sum of all distances between unordered pairs of distinct vertices of  $G$  [5]. The *average distance* of  $G$ , denoted by  $\bar{D} = \bar{D}(G)$ , is  $2W/[n(n-1)]$ . Put another way, this number gives, on average, the distance between a pair of vertices of  $G$ . Unless stated otherwise, when we refer to a subgraph of  $G$ , we mean an induced subgraph.

The total domination number of a graph was first introduced in [2]. This invariant remains of interest to researchers as evidenced by numerous recent papers. Various upper and lower bounds on  $\gamma_t$  have been discovered. The domination number has, of course, been well studied [8,9].

The average distance of a graph has sometimes been used to provide lower bounds for domination-related invariants, including the domination number itself [4]. One of the first results along these lines is the following theorem due to F. Chung in [1], which originated as a conjecture of the computer program Graffiti [6]. The *independence number* of  $G$ , denoted by  $\alpha(G)$ , is the cardinality of a largest set of mutually non-adjacent vertices.

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**Theorem 1** (Chung). *Let  $G$  be a graph. Then*

$$\geq \bar{D},$$

*with equality holding if and only if  $G$  is complete.*

Recently, this theorem has been generalized by Hansen et al. as a result about the *forest number*  $f = f(G)$  of a graph  $G$  [7]. This is the maximum order of an induced forest of  $G$ .

**Theorem 2** (Hansen et al.). *Let  $G$  be a graph. Then*

$$f \geq 2\bar{D}.$$

This theorem was also motivated by a conjecture of Graffiti [10]. Its proof is based on techniques introduced by Dankelmann in [3]. Dankelmann uses similar techniques in [4] to characterize graphs with fixed order and domination number that have maximum average distance. One can derive the following theorem as a corollary of this characterization (although this is not stated in [4]).

**Theorem 3.** *Let  $G$  be a graph. Then*

$$> \frac{2}{3}\bar{D} - \frac{1}{3}.$$

*Moreover, this inequality is best possible.*

The proof of Dankelmann's characterization result is lengthy and technical. We give a short direct proof of Theorem 3, as well as the following Theorem 4, which is the main result of our paper. We defer the proofs to a later section.

**Theorem 4** (Main Theorem). *Let  $G$  be a graph. Then*

$$t > \bar{D} - \frac{1}{2}.$$

*Moreover, this inequality is best possible.*

## 2. OTHER DEFINITIONS

Let  $R(k, t, l)$  denote the *binary star* on  $k + t + l$  vertices, where the maximal interior path has order  $t$  and there are  $k$  leaves on one side of the binary star and  $l$  leaves on the other. See Figure 1.

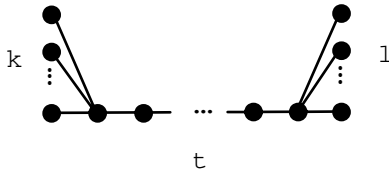


FIGURE 1. Binary star  $R(k, t, l)$ .

Now let  $R(n, t)$  denote the binary star of order  $n$  where the maximal interior path has order  $t$  and the leaves are as balanced as possible on each side of the binary star.

A set  $D$  of vertices of a graph  $G$  is called a *connected dominating set* provided  $D$  is a dominating set that induces a connected subgraph of  $G$ . The *connected*

*domination number* of  $G$ , denoted  $\gamma_c = \gamma_c(G)$ , is the cardinality of a smallest connected dominating set in  $G$ . A *trunk* for a graph  $G$  is a sub-tree (not necessarily induced) that contains the vertices of a dominating set of  $G$ . Hence, every spanning tree of  $G$  is a trunk for  $G$ , and every connected dominating set is the vertex set of some trunk. Standard graph theoretical terms not defined in this paper can be found in [11], for instance.

### 3. LEMMAS

The proof of Lemma 5 involves elementary algebra, counting, and limit arguments; we therefore omit it.

**Lemma 5.** *For integers  $k \geq 0$  and  $t \geq 1$ ,*

$$W(R(k, t, k)) = (t+3)k^2 + (t+2)(t-1)k + \frac{t(t+1)(t-1)}{6}, \text{ and}$$

$$W(R(k, t, k+1)) = (t+3)k^2 + (t+1)^2k + \frac{t(t+1)(t+2)}{6}.$$

Moreover,

$$W(R(k, t, k)) < W(R(k, t, k+1)) < W(R(k+1, t, k+1)), \text{ and}$$

$$\lim_{k \rightarrow \infty} \bar{D}(R(k, t, k)) = \frac{t+3}{2}.$$

The following lemma is proven in [6, Theorem 2].

**Lemma 6.** *Let  $G$  be a graph with a trunk of order  $t \geq 1$ . Then*

$$\bar{D}(G) \leq \bar{D}(R(n, t)),$$

*with equality holding if and only if  $G = R(n, t)$ .*

The next lemma follows by combining the two previous lemmas.

**Lemma 7.** *Let  $G$  be a graph with a trunk of order  $t \geq 1$ . Then*

$$\bar{D}(G) < \frac{t+3}{2}.$$

An immediate consequence of Lemmas 5 and 7 is the following corollary, which defines the relationship between the minimum order of a connected dominating set of a graph  $G$ , denoted  $\gamma_c = \gamma_c(G)$ , and its average distance.

**Corollary 8.** *Let  $G$  be a graph. Then*

$$\gamma_c > 2\bar{D} - 3.$$

Moreover, this inequality is best possible.

*Proof.* Let  $D$  be a minimum connected dominating set. Then any spanning tree of the subgraph induced by  $D$  is a trunk for  $G$ . Hence, by Lemma 7,

$$\bar{D}(G) < \frac{\gamma_c + 3}{2}.$$

To show this inequality is best possible, consider  $R(j, t, j)$ , where  $t \geq 1$  and  $j \geq 0$ . It is easy to see that  $\gamma_c(R(j, t, j)) = t$ . But by Lemma 5,

$$\lim_{j \rightarrow \infty} \bar{D}(R(j, t, j)) = \frac{t+3}{2} = \frac{\gamma_c + 3}{2}.$$

□

One final lemma is needed. The next simple lemma provides some relations that hold for the number of edges induced by dominating sets and their complements. Given a graph  $G$  with dominating set  $D$ , a vertex  $v \notin D$  is *over-dominated* by  $D$  if it has two or more neighbors in  $D$ . The *over-domination number of  $v$  with respect to  $D$* , denoted by  $O_D(v)$ , is one less than the number of neighbors  $v$  has in  $D$ .

**Lemma 9.** *Let  $T$  be a tree with minimum dominating set  $D$  such that the number of components of  $D$  is  $k$ . Denote the number of edges with both endpoints in  $D$  by  $e_1$ , the number of edges with both endpoints in  $H = T - D$  by  $e_2$ , and the number of edges with one endpoint in  $D$  and the other endpoint in  $H$  by  $e_3$ . Moreover, let  $j$  be the number of non-trivial components of  $H$  with at least two neighbors in  $D$  and let  $l_H$  be the number of components of  $H$  with exactly one neighbor in  $D$  (the leaves of  $H$ ). Then*

- a)  $e_1 = |D| - k$
- b)  $e_2 = k - 1 - \sum_{v \in H} O_D(v)$
- c)  $e_3 = n - |D| + \sum_{v \in H} O_D(v)$
- d)  $2j + l_H \leq e_3 = k + j + l_H - 1$
- e)  $n - l_H + 2 + \sum_{v \in H} O_D(v) \leq 2k + |D|$ .

*Proof.* Part a) holds because  $D$  induces a forest with  $k$  trees. Part c) is true because every vertex in  $H$  has a neighbor in  $D$ , giving the  $n - |D|$ , and because the summation contributes the extra edges that have one endpoint  $D$  and one in  $H$ . Part b) follows immediately from parts a) and c), since  $n - 1 = e_1 + e_2 + e_3$  for a tree.

The left hand side of d) comes from the fact that, when counting the edges between  $D$  and  $H$ , each of the  $l_H$  leaves in  $H$  contributes exactly one edge while each of the  $j$  non-trivial components of  $H$  contributes at least two edges. The right hand side of d) follows easily by viewing the components of  $D$  together with the components of  $H$  as the vertices of a new tree with  $e_3$  edges and  $k + j + l_H$  vertices.

From d) we deduce that there are at most  $k - 1$  non-trivial components of  $H$ , that is,  $j \leq k - 1$ . Combining this with the right hand side of d) and part c), we arrive at inequality e).  $\square$

#### 4. THEOREM PROOFS

Our strategy for proving Theorem 4 is as follows. Given a minimum total dominating set  $D$  of a graph  $G$ , we form a particular spanning tree  $T$  of  $G$  so that  $D$  is also a minimum total dominating set of  $T$ . Then we apply the lemmas from the previous section to obtain the desired result.

**Theorem 4 (Main Theorem)** *Let  $G$  be a graph. Then*

$$t > \bar{D} - \frac{1}{2}.$$

*Moreover, this inequality is best possible.*

*Proof.* Let  $D$  be a minimum total dominating set of  $G$ . Suppose that  $D$  has  $k$  components. We form a spanning tree  $T$  of  $G$  such that  $D$  is also a minimum total dominating set of  $T$ . If  $G$  is a tree, then put  $T = G$  and we are done. Otherwise, let  $C$  be a cycle in  $G$ . We delete an edge from  $C$  as follows.

- i) If  $C$  has two consecutive vertices  $x$  and  $y$  such that  $x \notin D$  and  $y \notin D$ , then delete the edge between them. The set  $D$  is still total dominating set for the resulting graph.
- ii) Suppose the first case does not apply. If  $C$  has two consecutive vertices  $x$  and  $y$  such that  $x \in D$  and  $y \notin D$ , then delete the edge between them. Since the other neighbor of  $y$  on  $C$  is necessarily in  $D$  (or else the first case applies), the set  $D$  is still a total dominating set for the resulting graph.
- iii) If neither of the first two cases apply, then all of the vertices of  $C$  are in  $D$ . Delete any edge of  $C$  and the set  $D$  is still a total dominating set for the resulting graph.

Repeat this process until all cycles are removed. Call the resulting spanning tree  $T$ . Since  $D$  is a total dominating set of  $T$ ,  $\tau(T) \leq |D| = \tau(G)$ . Since the total domination number of a graph is at most the total domination number of any of its spanning trees,  $\tau(G) \leq \tau(T)$ . Thus,  $\tau(T) = |D|$  and  $D$  is a minimum total dominating set of  $T$ .

Now, let  $L_H$ , of cardinality  $l_H$ , denote the leaves of  $T$  that are in  $H = T - D$  (the leaves of  $T$  that are not in  $D$ ). Observe that the sub-tree  $T - L_H$  contains the total dominating set  $D$  of  $G$  and is thereby a trunk for  $G$ . From Lemma 7,

$$2\bar{D} - 3 < |T - L_H| = n - l_H.$$

Hence by Lemma 9 part e), and since  $2k \leq \tau$ ,

$$2\bar{D} - 3 < 2k + \tau - 2 - \sum_{v \in H} O_D(v) \leq 2\tau - 2 - \sum_{v \in H} O_D(v) \leq 2\tau - 2.$$

Rearranging yields the desired inequality.

To show the inequality is best possible, consider  $R(j, t, j)$ , where  $t \equiv 2 \pmod{4}$  and  $j \geq 0$ . It is easy to see that  $\tau(R(j, t, j)) = \frac{t}{2} + 1$ . But by Lemma 5,

$$\lim_{j \rightarrow \infty} \bar{D}(R(j, t, j)) = \frac{t}{2} + \frac{3}{2} = \tau + \frac{1}{2}.$$

□

The proof of the theorem provides a necessary condition for  $\tau = \lceil \bar{D} - \frac{1}{2} \rceil$ . In the proof we found a spanning tree  $T$  of a connected graph  $G$  such that a minimum total dominating set of  $G$  was also a total dominating set for  $T$ . We let  $H = T - D$  and found that

$$\tau > \bar{D} - \frac{1}{2} + \frac{1}{2} \sum_{v \in H} O_D(v).$$

Now if  $\tau = \lceil \bar{D} - \frac{1}{2} \rceil$ , then

$$\lceil \bar{D} - \frac{1}{2} \rceil = \tau \geq \lceil \bar{D} - \frac{1}{2} + \frac{1}{2} \sum_{v \in H} O_D(v) \rceil,$$

which immediately suggests that  $D$  may over-dominate at most one vertex of  $H$ , and if there is an over-dominated vertex of  $H$ , its over-domination number is 1.

To see that there exist graphs in which any spanning tree containing a minimum total dominating set of the graph (as a total dominating set for the spanning tree) over-dominates exactly one vertex (with over-domination number 1) of  $H$  and  $\tau = \lceil \bar{D} - \frac{1}{2} \rceil$ , consider  $R(j, t, j)$ , where  $t > 1$ ,  $t \equiv 1 \pmod{4}$  and  $j \geq t$ . On the other

hand, that this condition is not sufficient for equality is seen in  $P_{4k+3}$  (the path on  $4k + 3$  vertices) for  $k \geq 1$ . Any minimum total dominating set  $D$  in  $P_{4k+3}$  over-dominates exactly one vertex  $v$  of  $V - D$ , and  $v$  has over-domination number 1, but  $\bar{d}$  is about one half the number of vertices and  $\bar{D}$  is about one third of the number vertices.

Next we present a short and direct proof of Theorem 3. As mentioned previously, this result can be deduced from a result of Dankelmann in [4].

**Theorem 3** *Let  $G$  be a graph. Then*

$$\bar{d} > \frac{2}{3}\bar{D} - \frac{1}{3}.$$

*Moreover, this inequality is best possible.*

*Proof.* Let  $D$  be a minimum dominating set of  $G$ . Suppose that  $D$  has  $k$  components. We will form a spanning tree  $T$  of  $G$  such that  $D$  is also a minimum dominating set of  $T$ . If  $G$  is a tree, then put  $T = G$  and we are done. Otherwise, let  $C$  be a cycle in  $G$ . We delete an edge from  $C$  as follows.

- i) If  $C$  has two consecutive vertices  $x$  and  $y$  such that  $x \notin D$  and  $y \notin D$ , then delete the edge between them. The set  $D$  still dominates the resulting graph.
- ii) Suppose the first case does not apply. If  $C$  has two consecutive vertices  $x$  and  $y$  such that  $x \in D$  and  $y \notin D$ , then delete the edge between them. Since the other neighbor of  $y$  on  $C$  is necessarily in  $D$  (or else the first case applies), the set  $D$  still dominates the resulting graph.
- iii) If neither of the first two cases apply, then all of the vertices of  $C$  are in  $D$ . Delete any edge of  $C$  and the set  $D$  still dominates the resulting graph.

Repeat this process until all cycles are removed. Call the resulting spanning tree  $T$ . Since  $D$  is a dominating set of  $T$ ,  $\bar{d}(T) \leq |D| = \bar{d}(G)$ . Since the domination number of a graph is at most the domination number of any of its spanning trees,  $\bar{d}(G) \leq \bar{d}(T)$ . Thus,  $\bar{d}(T) = |D|$  and  $D$  is a minimum dominating set of  $T$ .

Now, let  $L_H$ , of cardinality  $l_H$ , denote the leaves of  $T$  that are in  $H = T - D$  (the leaves of  $T$  that are not in  $D$ ). Observe that the sub-tree  $T - L_H$  contains the dominating set  $D$  of  $G$  and is thereby a trunk for  $G$ . From Lemma 7,

$$2\bar{D} - 3 < |T - L_H| = n - l_H.$$

Hence by Lemma 9 part e), and since  $2k \leq l_H$ ,

$$2\bar{D} - 3 < 2k + n - 2 - \sum_{v \in H} O_D(v) \leq 3n - 2 - \sum_{v \in H} O_D(v) \leq 3n - 2.$$

Rearranging yields the desired inequality.

To show the inequality is best possible, consider the family of stars  $S_n$ . Since the average distance in stars can be made arbitrarily close to 2,  $\frac{2}{3}\bar{D}(S_n) - \frac{1}{3}$  can be made arbitrarily close to  $\bar{d}(S_n) = 1$ .  $\square$

As was the case for total domination number and average distance, one can deduce from the proof a similar necessary condition for equality in  $\bar{d} = \lceil \frac{2}{3}\bar{D} - \frac{1}{3} \rceil$ .

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